# **Fast & Efficient Python Programming Workshop**

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## Why Python is slow ?

**Dynamic Typing in Python:**

**• Variables are dynamically typed**: They get their type at runtime when values (PyObject\*) are assigned.

- **Impact on Performance**:
	- This makes it difficult for the interpreter to optimize execution.
	- In contrast, compiled languages allow extensive analysis and optimization before runtime.



## Why Python is slow ?

#### **Dynamic Typing in Python:**

- Python has only one data type, PyObject\* with a pointer to its runtime type, which is yet another PyObject\*.
- Python is a dynamically typed language (Duck Typing). It wraps and later unwraps objects (referred to as boxing/unboxing).





## Why Python is slow ?

#### **Flexible Data Structures in Python**

- **Built-in Structures**: Python's built-ins (e.g., lists, dictionaries) are highly flexible and versatile.
- **Trade-offs**:
	- Generic nature, less efficient for numerical computations.
	- Perform well when processing diverse data types, but introduce significant overhead when handling large amounts of uniform data.



#### What to know, before optimising your code for acceleration …

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❏ How do we measure performance?

- ❏ Navigating algorithmic complexity
	- CPU Operations  $O(\log n)$  $O(1)$ Input Size
- ❏ Tools that help achieve this











#### **Why Execution Time Isn't the Best Metric for Algorithmic Complexity ?**

● **Execution time** alone is not a reliable measure of algorithmic complexity because it depends on external factors like **hardware**, **system load**, and **compiler optimizations**.

Objective complexity analysis requires a more standardized approach, such as evaluating the algorithm's asymptotic behavior (**Big O notation**), which focuses on **input size** rather than specific **execution time**.

● Instead of timing, we need to consider a **line-by-line evaluation** of the algorithm's structure, breaking down operations to assess how the number of steps grows as input size increases.



➔ **Line-by-line evaluation**

Python provides the ability to evaluate your code line by line, allowing you to identify **bottlenecks**.

● By pinpointing these bottlenecks, we can optimize our code, leading to increased **efficiency** and **faster performance**.

● Python offers several profiling libraries. We will focus on two key libraries: **cProfile** and **line\_profiler**.



➔ **Line-by-line evaluation**

Function which sorts a list of elements using the **bubble sort algorithm**.

```
import cProfile 
#Here we are using CProfile
def bubble_sort(a):
      n = len(a)for i in range(n):
            for j in range(n - i - 1):
                   if a[j] > a[j + 1]:
                         a[j], a[j + 1] = a[j + 1], a[j]return a
```
#### Total time: 15.256 seconds





➔ **Line-by-line evaluation**

Function which sorts a list of elements using the **bubble sort algorithm**.

```
from line_profiler import LineProfiler 
#Here we are using Line_Profile
```

```
def bubble_sort(a):
      n = len(a)for i in range(n):
             for j in range(n - i - 1):
                   if a[j] > a[j + 1]:
                          a[j], a[j + 1] = a[j + 1], a[j]return a
```




➔ **Line-by-line evaluation**

#### **Profiling Overhead**

- **CProfile**: This profiler operates with **lower overhead**, giving a general overview of function execution times. It allows for efficient profiling without significantly affecting the performance of the code being analyzed.
- **Line Profiler**: In contrast, this profiler incurs higher overhead because it tracks execution time for each individual line of code. As a result, the reported execution times may be **longer due to the additional processing required**.

#### **Measurement Focus**

- **CProfile**: This tool **measures the total execution time of function calls**, which **includes the time spent in any sub-functions**. Consequently, this can create the impression that functions are faster than they actually are since it aggregates all execution time into a single measurement.
- **Line Profiler**: This profiler concentrates on the **time taken by each line of code**, offering detailed insights into performance inefficiencies. It is particularly **useful for identifying bottlenecks within loops and complex sections** of code.

**CProfile** provides broad overview of function performance, **Line Profiler** delivers in-depth insights into execution time of each line of code.



#### ➔ **Big-O Notation**

Big-O notation describes the relationship between the size of the input to an algorithm and the number of steps required to execute it.

Unlike measuring performance for a specific instance (such as calculating  $fact(50)$ , Big-O focuses on how well an algorithm scales with:

- 1. **Increasing Input Size**: How the algorithm's performance changes as the input size grows.
- 2. **Type of Input**: How the algorithm's efficiency varies with different types of input data.

This approach provides a more comprehensive evaluation metric than assessing concrete execution time for a specific case.



➔ **Big-O Notation**

Simple Example: **Quick Sort**



At each level of recursion, we make comparisons and partition the array, which takes **O(n)** time. The recursion proceeds for approximately **log(n)** levels, because the a**rray size halves with each step**.

Therefore, the time complexity of Quicksort is **O(n log n)**.

def QuickSort(arr): if  $len(arr) \leq 1$ : return arr pivot =  $arr[len(arr)$  // 2]  $left = [x for x in arr if x < pi$ middle =  $[x \text{ for } x \text{ in } \text{arr if } x == \text{ pivot}]$ right =  $[x \text{ for } x \text{ in } \arctan x > \text{ pivot}]$ return QuickSort(left) + middle + QuickSort(right)

**Quick Sort Algorithm** 19 17 15 12 16 18 4 11 13

13

 $5 = 16$ 

15

 $\leq 13$ 

 $>=11$ 

 $12$ 

 $7 | 12 | 4$ 

 $\left| \downarrow \right|$  <=4

 $\leq 11$ 

 $-$  Pivot

 $>=16$ 

19 18

 $>=18$ 19

 $>=13$ 18 15 19 16



➔ **Big-O Notation**

Simple Example: **Merge Sort**

Merge Sort works by **dividing** the array into two halves. Each half is **recursively sorted**. The two sorted halves are then **merged** back together into a single sorted array. It's stable and guarantees O(n log n) time complexity but requires extra space for merging.



def merge\_sort(arr): if  $len(ar) > 1$ : mid =  $len(arr)$  // 2 left = arr[:mid]  $right = arr[mid.]$  # Recursively sort both halves merge\_sort(left) merge\_sort(right)  $i = j = k = 0$  # Merge the two halves while  $i <$  len(left) and  $j <$  len(right): if left[i] < right[j]:  $arr[k] = left[i]$  $i + 1$  else:  $arr[k] = right[j]$  $j + 1$  $k + = 1$  # Check if any element was left in the left half while i < len(left):  $arr[k] = left[i]$  $i + = 1$  $k + = 1$  # Check if any element was left in the right half while  $j$  < len(right):  $arr[k] = right[j]$  $j + 1$  $k + = 1$ 

- ➔ **Big-O Notation**
- **1. Memory Usage:**
	- **Quick Sort** is an **in-place sorting algorithm**, meaning it doesn't need extra memory for temporary arrays, while **Merge Sort** requires additional space for merging. This makes Quick Sort more efficient in terms of memory usage, especially for large datasets.

#### **2. Cache Efficiency:**

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**Quick Sort** has better **cache locality**, accessing memory sequentially and utilizing modern CPU cache more effectively. **Merge Sort** accesses memory in a scattered way, leading to more cache misses and slower performance.



# A bit about Numpy **NT NumPy**



- Stores **data in continuous memory blocks** for high performance.
- Allows operations on entire arrays/lists without explicit loops for increased speed (**Vectorization**).
- Supports operations on arrays of different sizes without manual size adjustment (**Broadcasting**).
- Provides fast, element-wise array operations (**ufuncs**).
- Integrates well with **Pandas** and **SciPy**.
- Able to perform complex mathematical computations such as linear algebra and Fourier transforms.
- Uses l**ess memory compared to traditional Python lists**, with precise control over data types.
- **Consistent** over different platforms and OS.
- Can be extended with C or Fortran for performance-critical tasks.











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### **Vectorization**

**Vectorization refers to the process of converting operations that typically process one element at a time, such as those in a loop, into operations that process multiple elements simultaneously.** 





- **Key aspects:** 
	- **Data parallelism** utilizing SIMD architecture (as implemented in GPUs).
	- Contiguous memory management.
- Benefits:

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- Performance improvements by reducing the number of interpreted loops in high-level languages.
- Lower memory footprint by minimizing temporary variable storage.
- Cleaner and more concise code, better readability.
- Utilized in tons of libraries (some we will discuss today).

#### Loop-wise versus Vectorization

#### # Loop-based calculation

```
# Create two arrays of size 1,000,000
a = np.random.randn(1 000 000)b = np.random.randn(1 000 000)# Loop-wise multiplication
def loop wise multiplication(a, b):
    result = np empty like(a) # Initialize an empty array with the same shape as 'a'
```
for  $i$  in range( $len(a)$ ):  $result[i] = a[i] * b[i]$ return result

# Time the loop-wise multiplication <sup>8</sup> timeit  $c = loop$  wise multiplication(a, b)

# Display the first 5 elements of the result  $c = loop$  wise multiplication(a, b) # Call the function to get the result print("First 5 elements of the result:", c[:5])

246 ms  $\pm$  11.4 ms per loop (mean  $\pm$  std. dev. of 7 runs, 1 loop each) First 5 elements of the result: [0.37046711 0.0320318 0.58237442 0.07603239 0.03643322]

#### # Vectorization on small set

```
# Create two arrays of size 10,000
a = np.random.randn(10000)b = np.random.randn(10000)
```
# Vectorized multiplication **Atimeth**  $c = a * b$  # This line is vectorized

# Display the first 5 elements of the result print("First 5 elements of the result:", c[:5]) #print(timeit.time())

 $4.02 \text{ }\mu\text{s} \pm 243 \text{ ns per loop}$  (mean  $\pm$  std. dev. of 7 runs, 100000 loops each) First 5 elements of the result: [0.26792758 0.61447971 0.01041451 0.07903875 0.05245871]

#### # Vectorization on larger set

# Create two arrays of size 1,000,000  $a = np.random.randn(1 000 000)$  $b = np.random.randn(1 000 000)$ 

# Vectorized multiplication "stime it  $c = a * b$  # This line is vectorized

# Display the first 5 elements of the result print("First 5 elements of the result:", c[:5]) #print(timeit.time())

1.39 ms  $\pm$  90.5 µs per loop (mean  $\pm$  std. dev. of 7 runs, 1000 loops each) First 5 elements of the result: [0.26792758 0.61447971 0.01041451 0.07903875 0.05245871]



### **Broadcasting**

**Broadcasting simplifies the handling of arrays with different dimensions by automatically 'broadcasting' the smaller array across the larger one so that they have compatible shapes.** # Define a 3x3 matrix matrix = np.array( $[1, 2, 3]$ .



- **Rules for broadcasting:**
	- To deem which two arrays are suitable for operations, NumPy compares the shape of the two arrays dimension-by-dimension starting from the trailing dimensions, working it's way forward. (from right to left)
	- Two dimensions are said to be compatible if **both of them are equal, or either one of them is 1.**
	- If both the dimensions are unequal and neither of them is 1, then NumPy will throw an error and halt.



## Auto Differentiation

**Automatic differentiation (AD or autograd) is a set of techniques to evaluate the derivative of a function with a mix of numerical and symbolic approaches.**

Every computer program executes a sequence of elementary arithmetic operations and elementary functions.

By **applying the chain rule repeatedly** to these operations, derivatives of arbitrary order can be computed automatically and efficiently.

Widely implemented in libraries - JAX (uses XLA to run on GPUs), integrated autograd systems in Tensorflow and PyTorch etc.

#### **How Autograd Works?**

- Autograd can be implemented using two main methods:
	- **Forward mode AD**: Computes derivatives from the input towards the output. Suitable when there are fewer inputs than outputs.
	- **Reverse mode AD** (often used in deep learning): Computes derivatives from the output back to the inputs. Suitable when there are fewer outputs than inputs, as in the case of a loss function in neural networks.
- **Computation Graph**: AD involves constructing a computation graph where **nodes represent operations or variable**s and **edges represent dependencies between these operations**. Forward pass computes the values, and the backward pass propagates derivatives.





#### **Autodifferentiation example using TensorFlow**

**Minimizing a Quadratic Function with TensorFlow** 



# Define a variable for our starting point  $x = tf.Variable(0.0, name='x')$ 

```
# Define the function f(x) = (x - 3)^2def f(x):return (x - 3) ** 2
```
optimizer =  $tf.$  optimizers. SGD(learning rate=0.1) # Stochastic Gradient Descent as optimizer

steps =  $50$ # steps or iterations for gradient descent

x values, f values =  $[]$ ,  $[]$  # Lists to store the values of x and  $f(x)$  for plotting

tf. Gradient Tape(): context manager records the operations performed on TensorFlow variables for automatic differentiation

```
# Perform gradient descent
for step in range(steps):
    with tf.GradientTape() as tape:
        # Monitor 'x'
        \text{tape-watch}(x)y = f(x)
```
# Get the gradient of f with respect to x  $grad = tape, gradient(y, x)$ 

```
# Update x to minimize f(x)optimizer.apply_gradients([(grad, x)])
# Store the values for visualization
x values.append(x.numpy())
f<sub>values.append(y.numpy())</sub>
```

```
# Print the current value of x and f(x)print(f"Step {step + 1}: x = {x.nump()}, f(x) = {y.nump()}'')
```
#### # Final output print(f"\nOptimal x found: {x.numpy()}")

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Optimal x found: 2.999957323074341



### Auto Differentiation - Comparison

Grad descent calculated over 1M steps. GPU starts to performs better as number of steps increases.



# Function to perform gradient descent with timing def optimize function (use apu=True): # Optionally place computation on the GPU  $device = 'GPU:0'$  if use apu else '/CPU:0' with tf.device(device): # Define a variable for our starting point  $x = tf.Variable(0.0, name='x')$ 

> # Define the function  $f(x) = (x - 3)^2$  $def f(x):$ return  $(x - 3)$  \*\* 2

# Set up the optimizer optimizer =  $tf.$ optimizers.SGD(learning rate=0.01)

# Number of steps for gradient descent  $steps = 1 000 000$ 

# Lists to store the values of x and  $f(x)$  for plotting  $x_values, f_values = []$ , []

# Timing start  $start_time = time.time()$ 

# Perform gradient descent for step in range(steps): with tf.GradientTape() as tape:  $\text{tape}.\text{watch(x)}$  $y = f(x)$ 

> $grad = tape.qradient(y, x)$ optimizer.apply\_gradients([(grad, x)])

# Store the values for visualization  $x$  values.append( $x$ .numpy()) f\_values.append(y.numpy())

# Timing end  $duration = time.time() - start time$ 

return x\_values, f\_values, duration